KILLER features of the BEAM

And what makes the BEAM a unique and powerful tool that really stands out!

Actor Model

Hundreds, Thousands, Millions of processes...

System limit can be between 1,024-134,217,727

They share **NOTHING**

Communicate through message passing

Demo 1 A hypothetical server

a simplified DEMO SERVER

```
accept_loop() ->
 receive
              %% to get a message in my mailbox
  Req -> dispatch(Req, self()) %% self() = my process id
 end,
 accept_loop().
accept(NumRequests)->
 collect responses(0, 0, 0, 0, NumRequests).
dispatch(Req, AcceptorPid) ->
  spawn(
  ?MODULE,
                      %% Module
  [Req, AcceptorPid]).
                  %% Arguments
```

a simplified DEMO SERVER

a simplified DEMO SERVER

```
handle request(Req, ParentPid) when is_integer(Req) ->
   Resp = count_to_1000_and_do_other_stuff_too(Req, 0),
   HandlerPid = self(),
   ParentPid ! {HandlerPid, Resp}
 end.
count_to_1000_and_do_other_stuff_too(_Req, 1000) -> ok;
count_to_1000_and_do_other_stuff_too(Req, C) ->
  case (Req rem 2) of
    0 -> binary:copy(<<Req/integer>>,300);
    1 \rightarrow binary:copy(<<(Req + 1)/integer>>,200)
  end,
  count_to_1000_and_do_other_stuff_too(Req, C+1).
```

Demo 2 server with a bug

a simplified DEMO SERVER

```
handle_request(Req, ParentPid) when is_integer(Req) ->
  case Req rem 100 of
    ∅ ->
       io:format("~n**** ~p [INF]*****n", [Req]),
       ParentPid ! dont_wait_for_me,
       handle_with_inf_loop_bug();
    Other ->
      Resp = count_to_1000_and_do_other_stuff_too(Req, 0),
      HandlerPid = self().
      ParentPid ! {HandlerPid, Resp}
  end.
 end.
handle_with_inf_loop_bug()->
   infinite loop(0).
infinite_loop(C) ->
 A = binary:copy(<<1>>,200),
 _B = math:sqrt(1235),
 infinite_loop(C+1).
```

BEAM Scheduler

Cooperative?

Cooperative at C level

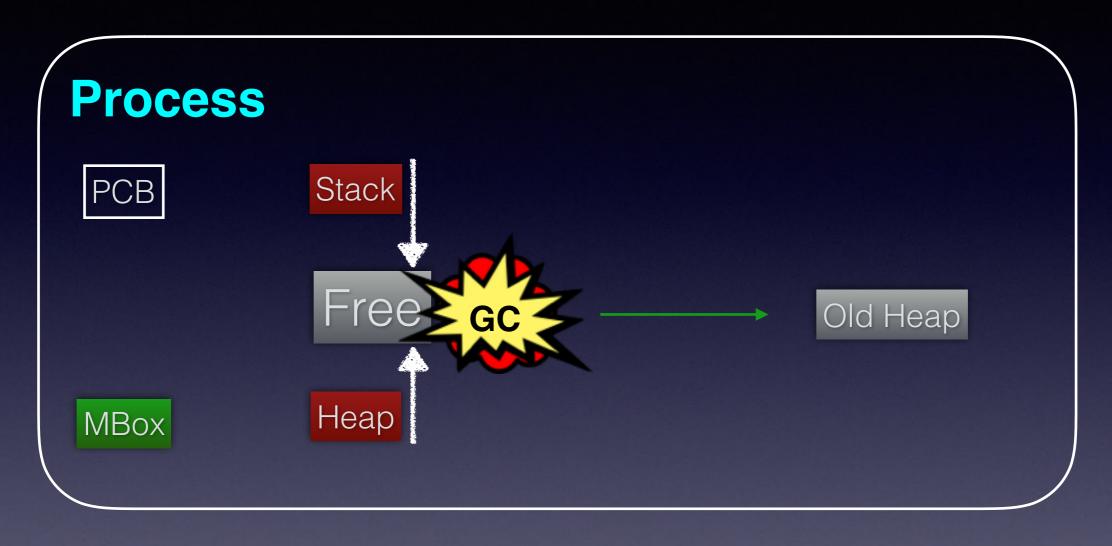
Preemptive?

Preemptive at Erlang level (by means of reduction counting)

2000 reductions
Reduction ~= function call

Word of caution: BIFs and NIFs

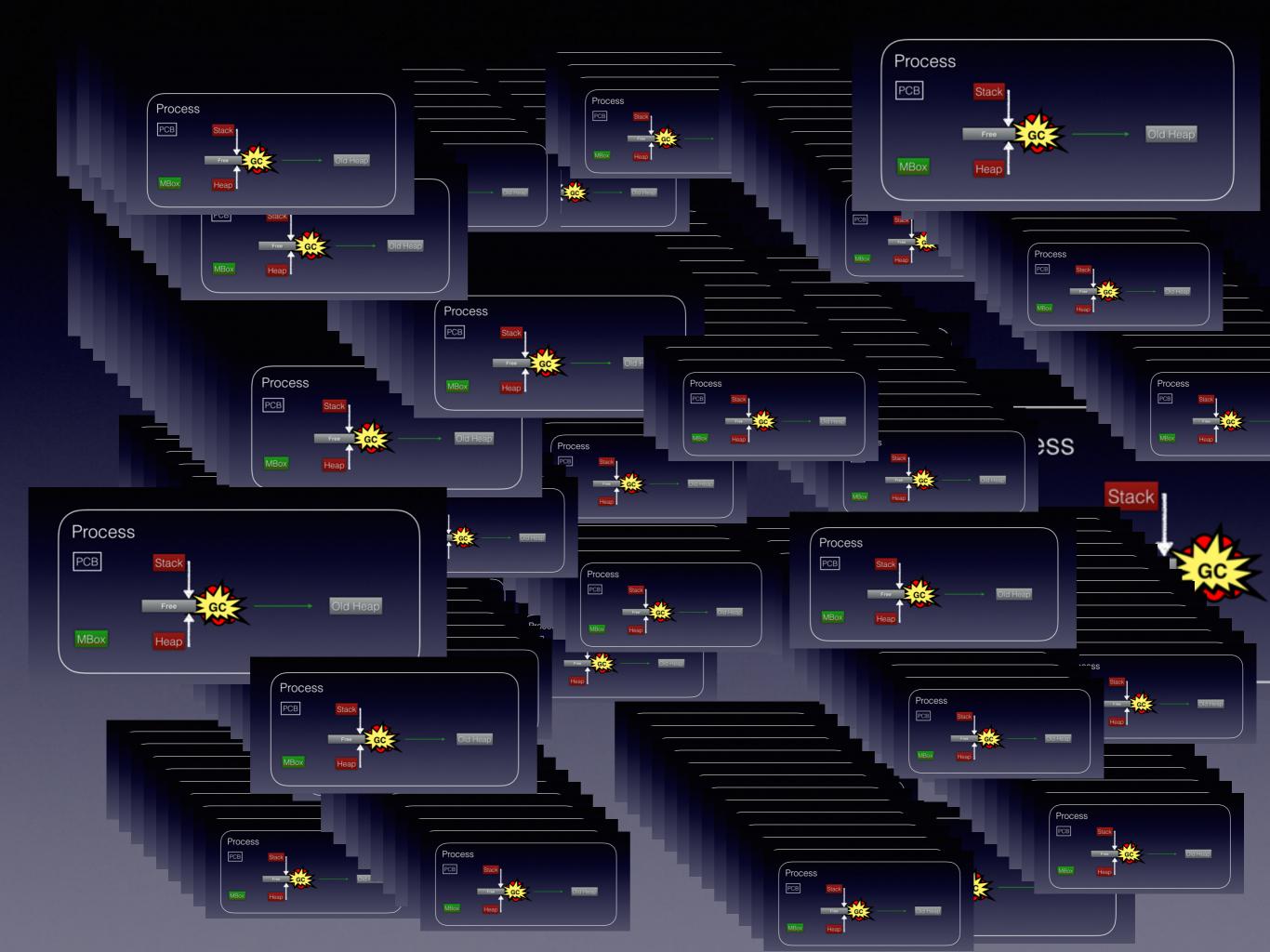
BEAM Memory model



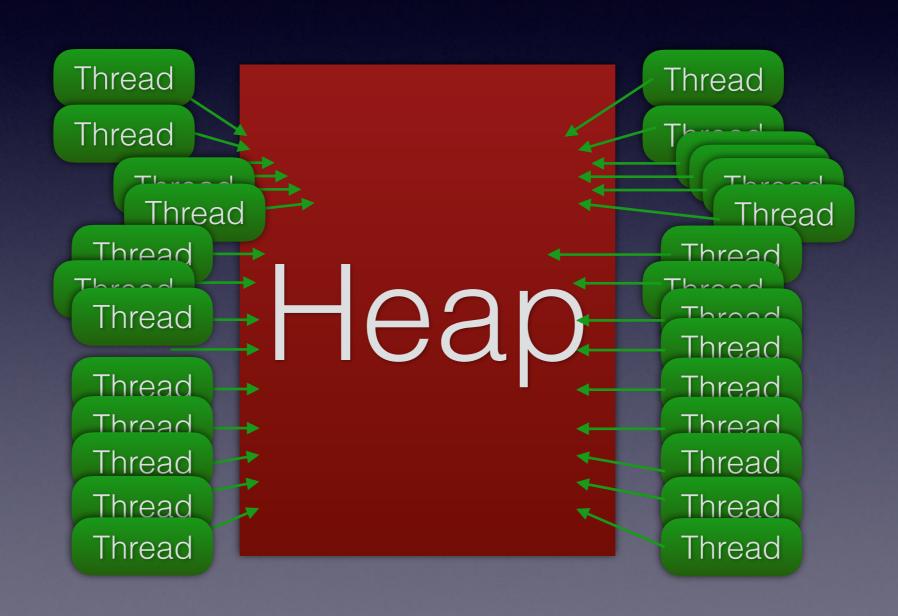
hipe_bifs:show_heap(Pid).

hipe_bifs:show_estack(Pid).

hipe_bifs:show_pcb(Pid). %% Look at heap_sz!



(simplified! SORRY!)





Stop the world. Is that a thing in Erlang?

DEMO3 The KILLER

KILLER_JUGGLER

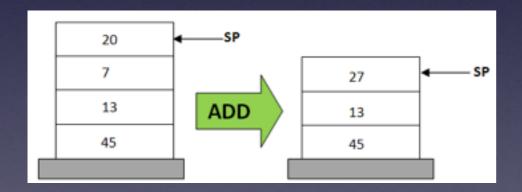
a simplified DEMO SERVER

```
%% a kind of a supervisor
kick_off_request_handler(Req, AcceptorPid) ->
  RequestHandlerPid =
    spawn(?MODULE, handle_request, [Req, self()]),
  Start = os:system_time(millisecond),
  receive
    {RequestHandlerPid, Resp} ->
         End = os:system_time(millisecond),
         Duration = End - Start,
         io:format("...~p [~b]...", [Req, Duration]),
        AcceptorPid ! {RequestHandlerPid, Resp, Duration};
    Other -> AcceptorPid ! {error, Other}
  after 5000 ->
    exit(HandlerPid, timedout),
    AcceptorPid ! {HandlerPid, killed}
  end.
```

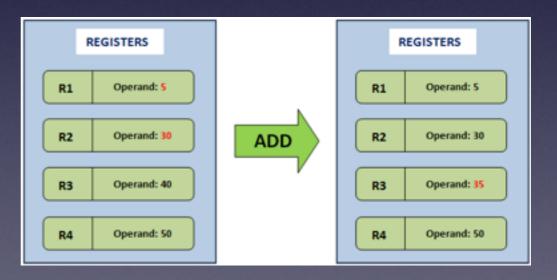
STACK BASED

REGISTER BASED

More Instructions Instructions are simple



1.POP 20 2.POP 7 3.ADD 20, 7, result 4.PUSH result Fewer instructions
Instructions have more info



1. ADD R1, R2, R3; # Add contents of R1 and R2, store result in R3

Performance Survey on Stack-based and Register-based VirtualMachines

Ruijie Fang, Siqi Liu, 2016

CONCEPTUM

(stack-based like JVM)

INERTIA

(register-based like BEAM)

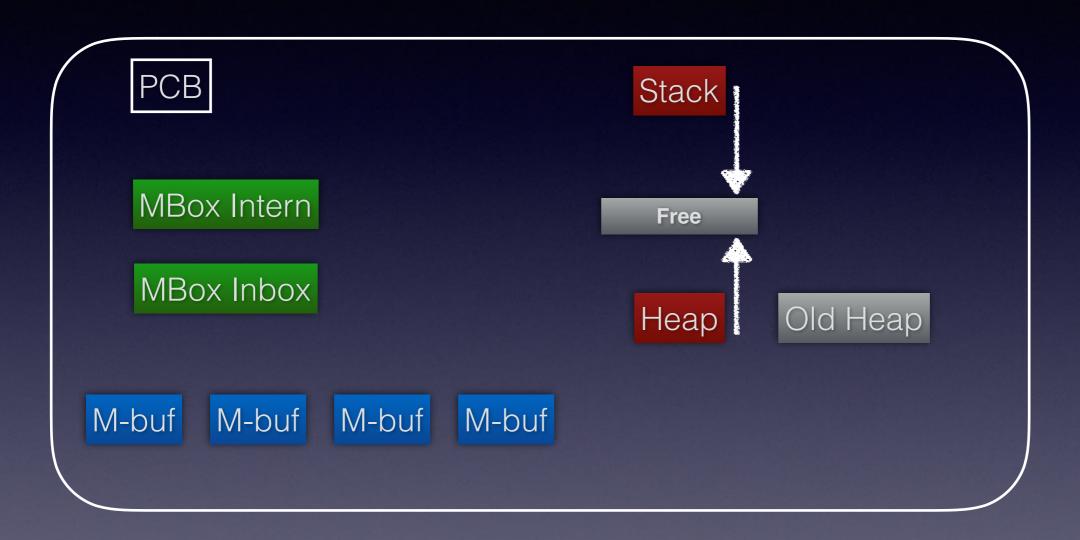
Inertia spends 66.42% less time in instruction dispatch than Conceptum, on average

However, Inertia is still slower in the overall fetch time, spending 23.5% more time on average in fetching operands than Conceptum does

Based on our test results, stack-based virtual machines typically perform better on benchmarks featuring a high amount of arithmetic operations.

In contrast to the stack-based virtual machine's performance, the register-based virtual machine performed much better on recursions and memory operations.

MBox

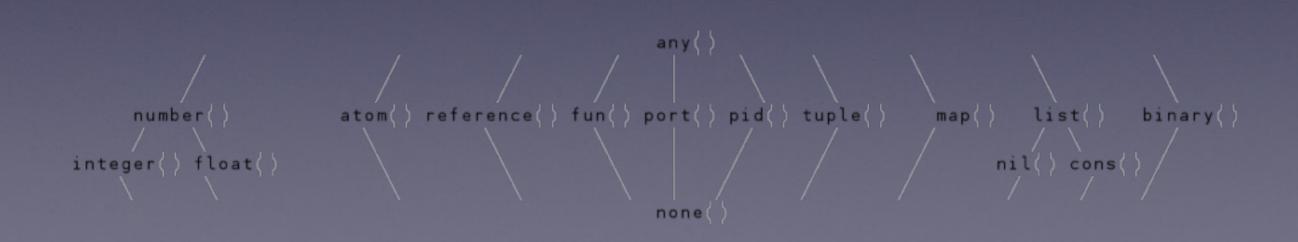


DEMO 4 Bring it down by sending messages to the process affected by inf loop!

TYPE SYSTEM

Strong typed.
So every type has a tag.

Dynamically typed. Hot code loading anyone?



TAGGING

In the memory representation of an Erlang term a few bits are reserved for a type tag.

```
LEVEL 1 tags:
  00 Header (on heap) CP (on stack)
  01 List (cons)
  10 Boxed <- pointers to the heap
  11 Immediate <- fit into one word on the stack
LEVEL 2 tags (Immediate):
  00 11 Pid
  01 11 Port
  10 11 Immediate 2
  11 11 Small integer
LEVEL 3 tags (Immediate 2):
 00 10 11 Atom
 01 10 11 Catch
 10 10 11 [UNUSED]
                             <- for empty list []
 11 10 11 Nil
```

Q. How is cooperative scheduling implemented?

A. If there are untagged values — no preempting

REFERENCES

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